

Probabilism, Entropies and Strictly Proper Scoring Rules

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Abstract

Accuracy arguments are the en vogue route in epistemic justifications of probabilism and further norms governing rational belief. These arguments often depend on the fact that the employed inaccuracy measure is strictly proper. I argue controversially that it is ill-advised to assume that the employed inaccuracy measures are strictly proper and that strictly proper statistical scoring rules are a more natural class of measures of inaccuracy. Building on work in belief elicitation I show how strictly proper statistical scoring rules can be used to give an epistemic justification of probabilism.

An agent's evidence does not play any role in these justifications of probabilism. Principles demanding the maximisation of a generalised entropy depend on the agent's evidence. In the second part of the paper I show how to simultaneously justify probabilism and such a principle. I also investigate scoring rules which have traditionally been linked with entropies.

Keywords: Accuracy, scoring rule, probabilism, strict propriety, entropy, principle of indifference

1 Introduction and Notation

2 1. Introduction

3 All Bayesians agree on one basic norm governing strength of rational
4 belief

5 **Probabilism:** Any rational agent’s subjective belief function
6 ought to satisfy the axioms of probability and every probability
7 function is, in principle, permissible. *Prob*

8 The question arises as to how to justify this norm. Traditionally, axiomatic
9 justifications [6, 41], justifications on logical grounds [22] and Dutch Book
10 Arguments [12, 50] were given to this end. Dutch Book Arguments have
11 been widely regarded as the most persuasive justification, however, they have
12 recently begun losing some of their once widespread appeal [21].¹

13 Recent epistemic justifications of probabilism are accuracy-based argu-
14 ments [24, 25, 30, 31, 49], which all build on [11]. The latter three arguments
15 employ Inaccuracy Measures (IMs) which are assumed to be strictly proper.
16 These IMs are closely related to the notion of a *Scoring Rule* (SR) which
17 the statistical community has a long tradition of studying, see [10] in the
18 Encyclopedia of Statistics.

19 In the first part of this paper, we argue that statistical SRs, properly
20 understood, are better suited than IMs to justify *Prob*. The argument will
21 be along the following lines: the most convincing justifications of *Prob* relying
22 on IMs require these IMs to be strictly proper (Section 4.1). However, for
23 the purposes of justifying *Prob*, assuming that an IM is strictly proper is ill-
24 advised (Section 4.3). On the contrary, assuming that a SR is strictly proper
25 is not only defensible but a desideratum (Section 3.2).

26 In Theorem 5.6 we show how strictly proper IMs give rise to strictly
27 proper SRs in a canonical way. We demonstrate in Theorem 6.2 how the
28 class of so-constructed SRs can be used to justify *Prob*.

¹We are joining the debate concerning rational belief formation assuming that degrees of beliefs are best represented by real numbers in the unit interval $[0, 1] \subset \mathbb{R}$. Anyone who rejects this premise will have to carefully assess whether the here presented account has implications on her line of thinking. Some of our results also hold true for degrees of belief represented by arbitrary positive real numbers.

29 The justifications in the first part of this paper do not take the agent's
 30 evidence into account. In all realistic cases rational agents do possess some
 31 evidence and this evidence ought to influence their degrees of belief, in some
 32 way. Maximum (generalised) entropy principles require an agent to adopt the
 33 probability function which maximises (a generalised) entropy among those
 34 probability functions which satisfy constraints imposed by her evidence.

35 In the second part of this paper we show how to simultaneously justify
 36 *Prob* and a such principle (Theorem 7.1 and Theorem 7.2). The usual argu-
 37 ment here consists of a two-stage justification – first one justifies *Prob* and
 38 then one justifies the entropy principle – and a story explaining why and how
 39 the justification of *Prob* trumps that of the entropy principle. The advantage
 40 of the simultaneous justification given here is that no such story needs to be
 41 told.

42 Taken together, *Prob* and such a principle entail the Principle of Indiffer-
 43 ence (PoI) in a large number of cases (Theorem 7.5, Corollary 7.6).

44 The logarithmic SR is well-known to be the only local SR which is strictly
 45 proper when applied to belief functions which are probability functions. Fur-
 46 thermore, this SR is at the heart of the maximum entropy principle. Since
 47 we here do not presuppose *Prob*, we investigate notions of locality applied
 48 to SRs for general belief functions (Section 8 and Section 9). We prove a
 49 non-existence result for such SRs in Theorem 8.4. Furthermore, we investi-
 50 gate how to weaken our assumptions to obtain strictly proper statistical SRs
 51 which are local in some sense, see Proposition 9.1 and Proposition 9.2.

52 2. The Formal Framework

53 Throughout, we work with a fixed, non-empty and finite set Ω , which is
 54 interpreted as the set possible worlds or elementary events. The power set
 55 of Ω , $\mathcal{P}\Omega$, is the set of events or the set of propositions. We shall assume
 56 throughout that $|\Omega| \geq 2$ and for $X \subseteq \Omega$ let $\bar{X} := \Omega \setminus X$.

57 The set of probability functions \mathbb{P} is the set of functions $P : \mathcal{P}\Omega \rightarrow [0, 1]$
 58 such that $\sum_{\omega \in \Omega} P(\{\omega\}) = 1$ and whenever $X \subseteq \Omega$ is such that $X = Y \cup Z$
 59 with $Y \cap Z = \emptyset$, then $P(X) = P(Y) + P(Z)$. We shall use $P(\omega)$ as shorthand
 60 for $P(\{\omega\})$.

61 Note that for all probability functions $P \in \mathbb{P}$ we have that $P(X) + P(\bar{X}) =$
 62 1 and hence $2 \sum_{X \subseteq \Omega} P(X) = \sum_{X \subseteq \Omega} P(X) + P(\bar{X}) = |\mathcal{P}\Omega|$.

63 The set of belief functions is the set of functions $Bel : \mathcal{P}\Omega \rightarrow [0, 1]$ and
 64 shall be denoted by \mathbb{B} . Throughout, we assume that all belief and probability

65 functions are *total*, i.e. defined on every $X \subseteq \Omega$. Trivially, since $|\Omega| \geq 2$ we
66 have $\mathbb{P} \subset \mathbb{B}$, where \subset denotes strict inclusion. Of particular interest are
67 the functions $v_\omega \in \mathbb{P}$ for $\omega \in \Omega$. A v_ω is the *at a world $\omega \in \Omega$ vindicated*
68 *credence function*. A v_ω can also be thought of as the *indicator function of*
69 *the elementary event $\omega \in \Omega$* . The v_ω are defined as follows:

$$v_\omega(X) := \begin{cases} 0 & \text{if } X \text{ is false at } \omega \\ 1 & \text{if } X \text{ is true at } \omega \end{cases} .$$

70 By “ X is true at ω ” we mean that $\omega \in X$; on the contrary, “ X is false at
71 ω ”, if and only if $\omega \notin X$.

72 In this paper we will stay within the classical framework of decision mak-
73 ing developed in [53]. So, we assume act-state independence², we also only
74 consider propositions which do not refer to themselves nor to their chances.
75 Such propositions are well-known to cause problems for the classical decision
76 making framework. Unsurprisingly, accuracy arguments based on the clas-
77 sical decision making framework are also troubled by such propositions, see
78 [5, 18]. Decision making frameworks for accuracy arguments which can deal
79 with such propositions are explored in [27].

80 Part 1

81 3. The Statistical Approach

82 3.1. Scoring Rules, Applications and Interpretations

83 Central to SRs and IMs is a measure function measuring the goodness or
84 badness, in some sense, of a belief function Bel . In the statistical community
85 this function is interpreted pragmatically as a loss incurred in a betting
86 scenario, whereas the epistemic tradition interprets the goodness measure as
87 a measure of (in)accuracy.

88 SRs have mainly been used to *elicit beliefs* or to *assess forecasts*. For
89 belief elicitation it is widely assumed that the agent’s belief function Bel^*

²In our context this means that neither the truth value nor the objective probability of a proposition $X \subseteq \Omega$ depends on the agent’s belief function Bel .

90 is a probability function, i.e., $Bel^* \in \mathbb{P}$. Similarly, forecasted events are
 91 normally assumed to be ruled by an objective probability function P^* , often
 92 taken to be the distribution of one (or several) random variable(s). In both
 93 applications, there exists a canonical probability function $P \in \mathbb{P}$ (either Bel^*
 94 or P^*) which can be used to aggregate losses incurred in different elementary
 95 events.

96 Formally, L is a loss function $L : \Omega \times \mathbb{P} \rightarrow [0, +\infty]$ and is referred to as a
 97 SR. For a guide to the voluminous literature to SRs refer to [17]. Expected
 98 loss is computed in the usual way

$$S_L : \mathbb{P} \times \mathbb{P} \rightarrow [0, +\infty], \quad S_L(P, Bel) := \sum_{\omega \in \Omega} P(\omega) \cdot L(\omega, Bel) . \quad (1)$$

99 Statisticians consider degrees of belief which satisfy *Prob*. Their notion
 100 of loss is thus only defined for probabilistic belief functions. For $Bel \in \mathbb{P}$ we
 101 have that Bel is completely determined by $\{Bel(\omega) \mid \omega \in \Omega\}$. In this case
 102 we can regard $L(\omega, Bel)$ as only depending on the first argument, ω , and
 103 $\{Bel(\omega) \mid \omega \in \Omega\}$.

104 We shall here be interested in *justifying Prob*. We thus consider a more
 105 general loss function L that also depends on degrees of belief in all non-
 106 elementary events $X \subseteq \Omega$. We thus consider a loss function $L : \Omega \times \mathbb{B} \rightarrow$
 107 $[0, +\infty]$ and define expected loss by

$$S_L : \mathbb{P} \times \mathbb{B} \rightarrow [0, +\infty], \quad S_L(P, Bel) := \sum_{\omega \in \Omega} P(\omega) \cdot L(\omega, Bel) . \quad (2)$$

108 In general, such a loss function $L : \Omega \times \mathbb{B} \rightarrow [0, +\infty]$ is *not* determined by
 109 the first argument, ω , and $\{Bel(\omega) \mid \omega \in \Omega\}$. Rather, $L(\omega, Bel)$ depends on
 110 the elementary event ω and $\{Bel(X) \mid X \subseteq \Omega\}$. So, although (1) and (2)
 111 appear at first glance to be the same expressions, they do differ in important
 112 aspects.

113 We shall tacitly assume that $L(\omega, Bel)$ in (1) and (2) may also depend
 114 on Ω throughout. That is, L may explicitly refer to the elementary events
 115 $\nu \in \Omega \setminus \{\omega\}$ or the the events $X \subseteq \Omega$ which contain ω . An example of the
 116 former kind of dependence can be found in (3) and of the latter kind in (14).

117 For ease of reading, we shall use the term *statistical SR* to refer to $S_L(\cdot, \cdot)$
 118 as in (2), rather than the long-winded “expectation of a SR $L : \Omega \times \mathbb{B} \rightarrow$

119 $[0, +\infty]$ ”.

120 The most famous SR is the Brier Score [3]:

121 **Definition 3.1.** *The Brier Score S_{Brier} takes the following form.³*

$$S_{Brier}(P, Bel) := \sum_{\omega \in \Omega} P(\omega) \cdot \left(\sum_{\mu \in \Omega} (v_\omega(\mu) - Bel(\mu))^2 \right) \quad (3)$$

$$= \sum_{\omega \in \Omega} P(\omega) \cdot \left((1 - Bel(\omega))^2 + \sum_{\mu \in \Omega \setminus \{\omega\}} Bel(\mu)^2 \right) \quad (4)$$

$$= \sum_{\omega \in \Omega} P(\omega) \cdot \left(1 - 2Bel(\omega) + \sum_{\mu \in \Omega} Bel(\mu)^2 \right) \quad (5)$$

$$= 1 + \sum_{\mu \in \Omega} Bel(\mu)^2 - \sum_{\omega \in \Omega} P(\omega) \cdot 2Bel(\omega). \quad (6)$$

122 See [57] for an axiomatic characterization of S_{Brier} .

123 3.2. Strict Propriety for statistical Scoring Rules

124 We now turn to the key property:

125 **Definition 3.2** (Strict \mathbb{X} -propriety). *For any set of belief functions $\mathbb{P} \subseteq \mathbb{X} \subseteq$*
 126 *\mathbb{B} , a statistical SR S_L is strictly \mathbb{X} -proper⁴, if and only if for all $P \in \mathbb{P}$*

$$\arg \inf_{Bel \in \mathbb{X}} S_L(P, Bel) = \{P\} . \quad (7)$$

127 In plain English, strictly \mathbb{X} -proper statistical SRs track probabilities,
 128 whatever these probabilities are.

³The original definition in [3] does not contain the formal expectation operator $\sum_{\omega \in \Omega} P(\omega) \cdot$. Rather, Brier envisioned a series of n forecasts which would all be scored by $\sum_{\omega \in \Omega} (Bel_i(\omega) - E_{i,\omega})^2$ where $Bel_i(\omega)$ notates the i -th forecast in ω and $E_{i,\omega}$ denotes indicator function for ω on the i -th occasion. The final score is then computed by dividing this sum by n . In essence, this amounts to taking expectations.

⁴Our notion of strict \mathbb{X} -propriety notably differs from Γ -strictness, see [20]. A SR is Γ -strict, if and only if for all $P \in \Gamma \subseteq \mathbb{P}$ it holds that $\arg \inf_{Bel \in \mathbb{P}} S_L(P, Bel) = \{P\}$; Γ -strictness is thus a weakening of strict \mathbb{P} -propriety. Strict \mathbb{B} -propriety is a strengthening of strict \mathbb{P} -propriety. $\Gamma \subseteq \mathbb{P}$ constraints the set of probability functions according to which expectations are computed, \mathbb{X} is a set of belief functions containing \mathbb{P} .

129 Recall from when we introduced statistical SRs that losses are usually
 130 interpreted pragmatically as losses in a betting scenario. For our purposes
 131 we will interpret the function S_L as a *measure of inaccuracy*. The intended
 132 interpretation is that $S_L(P, Bel)$ scores the inaccuracy of Bel with respect to
 133 the probability function P . By convention, score is an *inaccuracy* measure,
 134 a *low* score thus means low inaccuracy.

135 Now consider a function $P \in \mathbb{P}$ and a statistical SR $S_L(P, Bel)$. If
 136 $S_L(P, Bel)$ is strictly \mathbb{B} -proper, then $Bel = P$ is the unique belief func-
 137 tion for which $S_L(P, \cdot)$ is minimal. So, $Bel = P$ is the unique function
 138 which minimises inaccuracy. On the other hand, if $S_L(P, Bel)$ is not strictly
 139 \mathbb{B} -proper, then there exists a $P \in \mathbb{P}$ and a $Bel' \in \mathbb{B} \setminus \{P\}$ with $Bel' \in$
 140 $\arg \inf_{Bel \in \mathbb{B}} S_L(P, Bel)$. Arguably, then

141 The class of strictly \mathbb{B} -proper statistical SRs is the class of inac-
 142 curacy measures in the class of statistical SRs.

143 Plausibly, one might want to demand further desiderata (such as continuity
 144 of L) an inaccuracy measure ought to satisfy. However, it is not clear which
 145 other desideratum stands out in the class of further desiderata. Moreover, our
 146 approach covers the entire class of strictly \mathbb{B} -proper statistical SRs. We will
 147 henceforth take it that the class of statistical SRs which measure inaccuracy
 148 is the class of strictly \mathbb{B} -proper statistical SRs.

149 While S_{Brier} is well-known to be strictly \mathbb{P} -proper it is not strictly \mathbb{B} -
 150 proper since it does not depend at all on beliefs in non-elementary events
 151 and general belief functions $Bel \in \mathbb{B}$ are not determined by their values on
 152 elementary events. Thus, S_{Brier} cannot be the SR of choice for rational belief
 153 formation approaches that do not presuppose *Prob*.

154 To the best of our knowledge, strictly \mathbb{B} -proper SRs have, surprisingly, not
 155 been studied in the literature. So far, only strictly \mathbb{P} -proper SRs and strictly
 156 proper IMs (see Definition 4.2) have been investigated. In [29], Landes &
 157 Williamson use “strictly \mathbb{B} -proper SR” to refer to a function which computes
 158 expected losses of *normalised* belief functions. Their notion and our notion
 159 are thus not the same.

160 4. The Epistemic Approach

161 4.1. Ingredients

162 To highlight that we are now working within the epistemic framework we
 163 refer to the $\omega \in \Omega$ as possible worlds, Ω is now called the set of possible worlds

164 and the $X \subseteq \Omega$ are referred to as propositions. This change in terminology
 165 is, of course, purely cosmetic.

166 In recent epistemic approaches, the basic unit of inaccuracy is the inac-
 167 curacy of $Bel(X)$ at a world $\omega \in \Omega$, where proposition X is either true or
 168 false at ω . Formally, the inaccuracy is represented by an inaccuracy function
 169 $I(X, v_\omega(X), Bel(X))$. Since there may be reasons to treat different proposi-
 170 tions $X \subseteq \Omega$ differently, the inaccuracy of $Bel(X)$ at world ω may depend
 171 on the proposition $X \subseteq \Omega$. For example, different (additive or multiplica-
 172 tive) weights may be attached to different propositions. The basic inaccuracy
 173 units, $I(X, v_\omega(X), Bel(X))$, are then aggregated to an overall IM IM_I which
 174 measures the inaccuracy of $Bel \in \mathbb{B}$ with respect to a world $\omega \in \Omega$.

175 **Definition 4.1** (Inaccuracy Measure). *Let I be a function $I : \mathcal{P}\Omega \times \{0, 1\} \times$
 176 $[0, 1] \rightarrow [0, \infty]$. An IM IM_I is a map $IM_I : \Omega \times \mathbb{B} \rightarrow [0, \infty]$ such that*

$$IM_I(\omega, Bel) := \sum_{X \subseteq \Omega} I(X, v_\omega(X), Bel(X)) . \quad (8)$$

177 So, for a given world ω and a given belief function Bel , IM_I sums the
 178 inaccuracies over all propositions $X \subseteq \Omega$ of all beliefs $Bel(X)$ with respect to
 179 ω (or, depending on one’s point of view, with respect to the at ω vindicated
 180 credence function v_ω).

181 It is natural to think of I as some measure of distance between $v_\omega(X)$ and
 182 $Bel(X)$. For example, measuring inaccuracy in Euclidean terms one could
 183 consider

$$\begin{aligned} I(X, v_\omega(X), Bel(X)) &= (1 - Bel(X))^2, \text{ if } \omega \notin X \\ I(X, v_\omega(X), Bel(X)) &= Bel(X)^2, \text{ if } \omega \in X . \end{aligned}$$

184 Such an IM will formally be introduced in Definition 4.4.

185 The terminology in the literature has not yet converged. The function I
 186 has been called an (local) “inaccuracy measure” in [30, 43], whereas Predd
 187 et al. call I a SR and refer to IM_I as a “penalty function”, while Joyce calls
 188 it a “component function” in [25]. Groves (private communications) refers
 189 to I as “proposition-specific inaccuracy measure” which is more to the point
 190 but quite a mouthful.

191 In principle, it would be desirable to measure inaccuracy by some function

192 $f : \Omega \times \mathbb{B} \rightarrow [0, +\infty]$ (possibly satisfying further conditions) without assuming
 193 that f can be written as a sum over the $X \subseteq \Omega$. For further discussion on
 194 this point see [30, Section 5.2.1]. For the purposes of this paper we shall be
 195 interested in the set-up of Definition 4.1.

196 Conceptually, statistical SRs and IMs formalise notions of *inaccuracy*.
 197 While they share a common idea they measure inaccuracy differently. Statis-
 198 tical SRs measure inaccuracy between a belief function Bel and a probability
 199 function $P \in \mathbb{P}$, strictly \mathbb{B} -proper statistical SRs track probabilities. Whereas
 200 IMs measure inaccuracy between a belief function Bel and a possible world
 201 $\omega \in \Omega$, strictly proper IMs track the actual world, as we will see shortly. For
 202 some further discussion see Section 6.1.

203 One final difference of note is that $S_L(P, Bel)$ is a single real number,
 204 whereas $IM_I(\omega, Bel)$ is a tuple of real numbers, one real number for each
 205 $\omega \in \Omega$.

206 **Definition 4.2** (Strict Propriety). *An IM IM_I is called strictly proper, if*
 207 *and only if the following two conditions are satisfied*

- 208 • for all $p \in [0, 1]$ and all $\emptyset \subset X \subset \Omega$ it holds that $pI(X, 1, x) + (1 -$
 209 $p)I(X, 0, x)$ is uniquely minimized by $x = p$
- 210 • $I(\Omega, 1, x) + I(\emptyset, 0, y)$ is uniquely minimised by $x = 1$ and $y = 0$.

211 Intuitively, strict propriety ensures that setting degrees of belief in X
 212 equal to the probability of X is the only way to minimise expected inaccuracy,
 213 see further Section 4.3.

214 In general, the second condition above is required because $P(\emptyset) = 0$ and
 215 $P(\Omega) = 1$ for all $P \in \mathbb{P}$ and later on we want p to equal the probability of X .

216 Some authors do not allow I to depend on X , see for instance [44]. For
 217 such a loss function the requirement that $I(1, x) + I(0, y)$ is uniquely min-
 218 imised by $x = 1$ and $y = 0$ is simply an instance of the first condition. For
 219 such an I , the second condition follows from the first.

220 If IM_I is strictly proper, then for all $\omega \in \Omega$ and all $X \subseteq \Omega$ such that
 221 $\omega \in X$ it holds that $I(X, 1, Bel(X)) + I(\bar{X}, 0, Bel(\bar{X}))$ is minimised, if and
 222 only if $Bel(X) = 1$ and $Bel(\bar{X}) = 0$. That is, Bel and v_ω agree on X and \bar{X} .
 223 Hence, $IM_I(\omega, Bel)$ is uniquely minimized by $Bel = v_\omega$. So, if $\omega^* \in \Omega$ is the
 224 actual world, then the strictly least inaccurate belief function is $Bel = v_{\omega^*}$.
 225 In this sense, strictly proper IMs track the actual world.

226 Strict propriety as a desideratum for IMs has been argued for in various
 227 contexts in which *Prob* is pre-supposed, see [14, 16, 19, 38]. We shall not

228 advance arguments for strict propriety here; in Section 4.3 we shall argue
 229 *against* the use of strictly proper IMs in the current context.

230 The following condition strikes us as a sensible property an IM should
 231 satisfy:

232 **Definition 4.3.** *An IM IM_I is called continuous, if and only if I is contin-*
 233 *uous in $Bel(X)$.*

234 Continuity is here taken in the usual sense: For all $X \subseteq \Omega$, for all $i \in$
 235 $\{0, 1\}$ and for all sequences $(Bel_n(X))_{n \in \mathbb{N}}$ converging to $Bel(X) \in [0, 1]$ it
 236 holds that $\lim_{n \rightarrow \infty} I(X, i, Bel_n(X)) = I(X, i, Bel(X))$, where both sides of
 237 this equation may be equal to $+\infty$.

238 The most popular IM is an epistemic version of the Brier Score S_{Brier} :

239 **Definition 4.4** (Brier IM). *The Brier IM is defined as*

$$IM_{Brier}(\omega, Bel) := \sum_{X \subseteq \Omega} (v_\omega(X) - Bel(X))^2 . \quad (9)$$

240 In other words: $IM_{Brier}(\omega, Bel)$ is the square of the Euclidean distance in
 241 $\mathbb{R}^{|\mathcal{P}\Omega|}$ between v_ω and Bel . It is well-known that IM_{Brier} is strictly proper and
 242 continuous. Recently, quadratic IMs, such as IM_{Brier} , have been advocated
 243 in [30, 31] on the grounds that they are the only class of measures which keep
 244 an agent out of certain epistemic dilemmas.

245 Compare this measure IM_{Brier} to S_{Brier} (Definition 3.1) and observe that
 246 $IM_{Brier}(\omega, Bel)$ depends on the entire belief function while $S_{Brier}(P, Bel)$
 247 only depends on beliefs in elementary events. In Definition 5.2, we will see
 248 how to associate IM_{Brier} and a statistical SR. For now, we simply observe
 249 the following structural similarity

$$S_{Brier}(v_\omega, Bel) = \sum_{\mu \in \Omega} (v_\omega(\mu) - Bel(\mu))^2$$

$$IM_{Brier}(\omega, Bel) = \sum_{X \subseteq \Omega} (v_\omega(X) - Bel(X))^2 .$$

250 4.2. Justifications of Probabilism

251 In justifications of norms of rational belief formation employing IMs it is
 252 normally assumed that the agent has no information as to which world is the

253 actual one. How is one then to aggregate inaccuracies $IM_I(\omega, Bel)$ in differ-
 254 ent worlds? Surely, one could simply add the inaccuracies up, $\sum_{\omega \in \Omega} IM_I(\omega, Bel)$.
 255 But why should one not multiply the inaccuracies, $\prod_{\omega \in \Omega} IM_I(\omega, Bel)$, or con-
 256 sider the sum of the logarithms of the inaccuracies, $\sum_{\omega \in \Omega} \log(IM_I(\omega, Bel))$?
 257 Apparently, there is no canonical way to aggregate the inaccuracies $IM_I(\omega, Bel)$
 258 for the possible worlds $\omega \in \Omega$.

259 The Decision Theoretic Norm (DTN) which is widely applied in such a
 260 situation is dominance. Historically, the first justification of *Prob* applying
 261 dominance was:

262 **Theorem 4.5** (De Finetti [11]).

- 263 • For all $Bel \in \mathbb{B} \setminus \mathbb{P}$ there exists some $P \in \mathbb{P}$ such that for all $\omega \in \Omega$
 264 $IM_{Brier}(\omega, Bel) > IM_{Brier}(\omega, P)$.
- 265 • For all $Bel \in \mathbb{P}$ and all $Bel' \in \mathbb{B} \setminus \{Bel\}$ there exists an $\omega \in \Omega$ such
 266 that $IM_{Brier}(\omega, Bel') > IM_{Brier}(\omega, Bel)$.

267 De Finetti's result relies on IM_{Brier} to measure inaccuracy. Plausibly,
 268 there are other IMs which measure inaccuracy. Recently, the following gen-
 269 eralisation has been proved in the context of belief *elicitation*:

270 **Theorem 4.6** (Predd et al. [49]). *If IM_I is a continuous and strictly proper*
 271 *IM, then:*

- 272 • For all $Bel \in \mathbb{B} \setminus \mathbb{P}$ there exists some $P \in \mathbb{P}$ such that for all $\omega \in \Omega$
 273 $IM_I(\omega, Bel) > IM_I(\omega, P)$.
- 274 • For all $Bel \in \mathbb{P}$ and all $Bel' \in \mathbb{B} \setminus \{Bel\}$ there exists an $\omega \in \Omega$ such
 275 that $IM_I(\omega, Bel') > IM_I(\omega, Bel)$.

276 Predd et al. credit Lindley (see [34]) for a precursor of their result.

277 The first parts of these theorems say that every non-probabilistic belief
 278 function $Bel \in \mathbb{B} \setminus \mathbb{P}$ is strongly accuracy dominated by some probability
 279 function and thus impermissible. The second parts mean that every proba-
 280 bilistic belief function $Bel \in \mathbb{P}$ is permissible, because no $Bel \in \mathbb{P}$ is weakly
 281 accuracy dominated.

282 The two other main justifications of *Prob* along similar lines are due to
 283 Joyce, see [24] and [25]. Both justifications apply dominance as DTN in the
 284 same way as de Finetti and Predd et al.

285 The former justification in [24], does not require that a measure of in-
286 accuracy $f(\omega, Bel)$ can be written as a sum over the propositions $X \subseteq \Omega$.
287 In order to prove the theorem Joyce has to assume a number of properties
288 f has to satisfy. The assumed symmetry property has been objected to in
289 [16, 35], Maher also objected to the convexity property. In his 2009 paper,
290 Joyce concedes that the objections raised have merit and that it would be
291 best to do without these properties [25, p. 285].

292 The latter justification ([25, Theorem 2]) also does not require that the
293 measure of inaccuracy $f(\omega, Bel)$ can be written as a sum over the propositions
294 $X \subseteq \Omega$. It is only assumed that the measure of inaccuracy f satisfies a
295 number of conditions one of which is that f has to be finitely-valued.

296 We feel that the main draw-back with [25, Theorem 2] is that it only
297 applies for every partition of propositions and not to all propositions $X \subseteq \Omega$.
298 For further discussions see [61, Section 1].

299 4.3. *Strict Propriety for Justifications of Probabilism*

300 We now argue that Theorem 4.6 does not provide a satisfactory justifi-
301 cation of *Prob* for belief formation. The problem lies with the requirement
302 that IM_I be strictly proper.

303 We fully agree with Joyce

304 [...] we cannot hope to justify probabilism by assuming that ratio-
305 nal agents should maximize the expected accuracy of their opin-
306 ions because the concept of an expectation really only makes sense
307 for agents whose partial beliefs already obey the laws of proba-
308 bility. [24, p. 590]

309 Proponents of strictly proper IMs may object that strict propriety guaran-
310 tees that it is permissible to hold degrees of belief that agree with known
311 probabilities.

312 This objection misses the mark in at least two decisive ways.

313 Firstly, a function f ought to be considered as a measure of inaccuracy in
314 virtue of f measuring inaccuracy and emphatically not solely on the virtue of
315 the belief functions it renders permissible given a certain DTN. This objection
316 does not make clear why every appropriate measure of inaccuracy IM_I has
317 to be strictly proper. Intuitively plausible properties such as $I(X, 1, x)$ has
318 a unique minimum on $[0, 1]$ for $x = 1$ or that $I(X, 1, x)$ is a (strictly) decreasing
319 function in $x \in [0, 1]$ do not feature in this objection.

320 Secondly, as Joyce already pointed out, why would an agent with a non-
 321 probabilistic belief function $Bel^* \in \mathbb{B} \setminus \mathbb{P}$ care for the following expectation
 322 $Bel^*(X)I(X, 1, Bel^*(X)) + (1 - Bel^*(X))I(X, 0, Bel^*(X))$? It seems that
 323 such an agent rather cares for the “expectation” $Bel^*(X)I(X, 1, Bel^*(X)) +$
 324 $Bel^*(\bar{X})I(X, 0, Bel^*(X))$. Since we are in the business of *justifying Prob*,
 325 an agent with degrees of belief $Bel^*(X) = 0$ for all $X \subseteq \Omega$ would not be
 326 threatened in her beliefs by strict propriety.

327 We conclude that assuming strict propriety for our purposes is ill-advised.
 328 So, Theorem 4.6 does *not* yield a satisfactory justification of *Prob* for belief
 329 *formation*.

330 4.4. Strict Propriety for Belief Elicitation

331 In the belief *elicitation* framework of Predd et al. it is assumed that the
 332 agent’s belief function Bel^* is a probability function. Predd et al. [49, p.
 333 4786] motivate strict propriety by “Our scoring rule thus encourages sincer-
 334 ity since your interest lies in announcing probabilities that conform to your
 335 beliefs.” That is, a subjective Bayesian agent avoiding inaccurate beliefs
 336 has a clear impetus to minimise the expectation $Bel^*(X)I(X, 1, Bel'(X)) +$
 337 $Bel^*(\bar{X})I(X, 0, Bel'(X))$ by announcing $Bel'(X) = Bel^*(X)$. I hence find
 338 no fault with the requirement of “strict propriety” for eliciting beliefs from
 339 subjective Bayesian agents, although I do object to it for the purposes belief
 340 formation.

341 Belief elicitation is at heart an empirical problem, which is often tackled
 342 by employing questionnaires, by conducting interviews and/or by observa-
 343 tional studies (of subjects playing [incentive compatible] games). SRs have
 344 made their way into the applied sciences [39, 65]. See [16, Section 3] for a
 345 recent philosophical treatment of belief elicitation.

346 5. Associating Inaccuracy Measures with Scoring Rules

347 5.1. Extended Scoring Rules

348 In this section we shall introduce a class of statistical SRs which allow us
 349 to connect IMs to the here introduced class of statistical SRs. We follow [29]
 350 and define:

351 **Definition 5.1** (Extended Scoring Rule). *A statistical SR $S_L : \mathbb{P} \times \mathbb{B} \rightarrow$*
 352 *$[0, \infty]$ is called extended, if and only if it can be written as*

$$S_L^{ext}(P, Bel) = \sum_{\omega \in \Omega} P(\omega) \cdot L(\omega, Bel) \quad (10)$$

$$= \sum_{X \subseteq \Omega} P(X) \cdot L'(X, Bel) \quad (11)$$

$$= \sum_{\omega \in \Omega} P(\omega) \cdot \sum_{\substack{X \subseteq \Omega \\ \omega \in X}} L'(X, Bel) \ , \quad (12)$$

353 for some function $L' : \mathcal{P}\Omega \times \mathbb{B} \rightarrow [0, \infty]$.

354 The name *extended* is somewhat unfortunate. Originally, it was intended
 355 to capture the fact that the domain of the SR has been *extended* from $\mathbb{P} \times \mathbb{P}$
 356 to $\mathbb{P} \times \mathbb{B}$ and that the sum in (10) is over all events $X \subseteq \Omega$ and not merely
 357 over the elementary events $\omega \in \Omega$ as in (1).

358 For our running example, Brier Scores, we give the following extended
 359 SR:

Definition 5.2 (Extended Brier Score).

$$S_{Brier}^{ext}(P, Bel) := \sum_{X \subseteq \Omega} P(X) \cdot \left((1 - Bel(X))^2 + Bel(\bar{X})^2 \right) \quad (13)$$

$$= \sum_{\omega \in \Omega} P(\omega) \cdot \left(\sum_{\substack{X \subseteq \Omega \\ \omega \in X}} (1 - Bel(X))^2 + \sum_{\substack{Y \subseteq \Omega \\ \omega \notin Y}} Bel(Y)^2 \right) \quad (14)$$

$$= \sum_{\omega \in \Omega} P(\omega) \cdot IM_{Brier}(\omega, Bel) \ . \quad (15)$$

360 **Proposition 5.3.** S_{Brier}^{ext} is strictly \mathbb{B} -proper.

361 *Proof.* The idea is to decompose $S_{Brier}^{ext}(P, Bel)$ into pairs of summands,
 362 where each pair is of the form $P(X) \cdot ((1 - Bel(X))^2 + Bel(\bar{X})^2) + P(\bar{X}) \cdot$
 363 $((1 - Bel(\bar{X}))^2 + Bel(X)^2)$. We then show that each such pair is uniquely
 364 minimised by $Bel(X) = P(X)$ and $Bel(\bar{X}) = 1 - P(X)$.

365 Consider the following minimization problem for fixed $P \in \mathbb{P}$, fixed $X \subseteq$
 366 Ω and $x := Bel(X)$, $y := Bel(\bar{X})$

$$\text{minimize} \quad P(X) \cdot ((1 - x)^2 + y^2) + (1 - P(X)) \cdot ((1 - y)^2 + x^2)$$

subject to $x, y \in [0, 1]$.

367 Note that the objective function of this minimisation problem is equal to
 368 $x^2 - 2xP(X) + P(X) + y^2 - 2y(1 - P(X)) + (1 - P(X))$. The unique minimum
 369 obtains for $x = P(X)$ and $y = 1 - P(X)$.

370 Hence, $Bel = P$ uniquely minimizes $S_{Brier}^{ext}(P, \cdot)$. \square

371 A version of de Finetti's Theorem (Theorem 4.5) for S_{Brier}^{ext} follows as a
 372 simple Corollary:

373 **Corollary 5.4.**

374 • For all $Bel \in \mathbb{B} \setminus \mathbb{P}$ there exists some $P \in \mathbb{P}$ such that for all $Q \in \mathbb{P}$
 375 $S_{Brier}^{ext}(Q, Bel) > S_{Brier}^{ext}(Q, P)$.

376 • For all $Bel \in \mathbb{P}$ and all $Bel' \in \mathbb{B} \setminus \{Bel\}$ there exists a $P \in \mathbb{P}$ such that
 377 $S_{Brier}^{ext}(P, Bel') > S_{Brier}^{ext}(P, Bel)$.

378 *Proof.* 1) Let $Bel \in \mathbb{B} \setminus \mathbb{P}$. By Theorem 4.5 there exists a $P_{Bel} \in \mathbb{P}$ such that
 379 for all $\omega \in \Omega$ it holds that $IM_{Brier}(\omega, Bel) > IM_{Brier}(\omega, P_{Bel})$. Using (15),
 380 the fact that Ω is finite and that for all $Q \in \mathbb{P}$ there exists an $\omega \in \Omega$ with
 381 $Q(\omega) > 0$ we find that $S_{Brier}^{ext}(Q, Bel) > S_{Brier}^{ext}(Q, P_{Bel})$.

382 2) We saw in Proposition 5.3 that S_{Brier}^{ext} is strictly \mathbb{B} -proper. Hence,
 383 $S_{Brier}^{ext}(Bel, \cdot)$ is uniquely minimised by $Bel = Bel$. \square

384 Note that de Finetti's Theorem applies dominance with respect to the
 385 possible worlds $\omega \in \Omega$ while the above corollary applies dominance with
 386 respect to the probability functions $Q \in \mathbb{P}$.

387 5.2. The Canonical Association

388 In this section we shall see how to canonically associate with every IM an
 389 extended SR. We shall give two further examples to illustrate the association.

390 **Definition 5.5** (Canonical Association). For IM_I define an associated sta-
 391 tistical SR S_I^{aso} by:

$$S_I^{aso}(P, Bel) := \sum_{\omega \in \Omega} P(\omega) \cdot IM_I(\omega, Bel) \quad (16)$$

$$= \sum_{\omega \in \Omega} P(\omega) \cdot \left(\sum_{\substack{X \subseteq \Omega \\ \omega \in X}} I(X, 1, Bel(X)) + \sum_{\substack{Y \subseteq \Omega \\ \omega \notin Y}} I(Y, 0, Bel(Y)) \right) \quad (17)$$

$$= \sum_{X \subseteq \Omega} P(X) \cdot I(X, 1, Bel(X)) + P(\bar{X}) \cdot I(X, 0, Bel(X)) \quad (18)$$

$$= \sum_{X \subseteq \Omega} P(X) \cdot \left(I(X, 1, Bel(X)) + I(\bar{X}, 0, Bel(\bar{X})) \right) . \quad (19)$$

392 So, letting $L'(X, Bel) := I(X, 1, Bel(X)) + I(\bar{X}, 0, Bel(\bar{X}))$ we see that
 393 S_I^{aso} is an extended SR.

394 For a fixed IM IM_I , $S_I^{aso}(P, Bel)$ is simply the expected inaccuracy of Bel ,
 395 where expectations are computed with respect to the probability function
 396 $P \in \mathbb{P}$.

397 **Theorem 5.6.** *IM_I is strictly proper, if and only if S_I^{aso} is strictly \mathbb{B} -proper.*

398 *Proof.* If IM_I is strictly proper, then for every $\emptyset \subset X \subset \Omega$ and all $P \in \mathbb{P}$

$$P(X) \cdot I(X, 1, Bel(X)) + P(\bar{X}) \cdot I(X, 0, Bel(X))$$

399 is uniquely minimised by $Bel(X) = P(X)$.

400 Furthermore, $I(\Omega, 1, Bel(\Omega)) + I(\emptyset, 0, Bel(\emptyset))$ is uniquely minimised by
 401 $Bel(\Omega) = 1$ and $Bel(\emptyset) = 0$. Applying (18) we now find that $S_I^{aso}(P, \cdot)$ is
 402 uniquely minimised by $Bel = P$.

403 Now, suppose that S_I^{aso} is strictly \mathbb{B} -proper. Then for all $p \in [0, 1]$ and
 404 all $P \in \mathbb{P}$ with $P(\omega) = p$ and $P(\omega') = 1 - p$ for different $\omega, \omega' \in \Omega$ we have

$$\begin{aligned} S_I^{aso}(P, Bel) &= \sum_{X \subseteq \Omega} P(X) \cdot I(X, 1, Bel(X)) + P(\bar{X}) \cdot I(X, 0, Bel(X)) \\ &= \sum_{\substack{U \subseteq \Omega \\ \omega, \omega' \in U}} 1 \cdot I(U, 1, Bel(U)) + 0 \cdot I(U, 0, Bel(U)) \\ &\quad + \sum_{\substack{W \subseteq \Omega \\ \omega, \omega' \notin W}} 0 \cdot I(W, 1, Bel(W)) + 1 \cdot I(W, 0, Bel(W)) \\ &\quad + \sum_{\substack{Y \subseteq \Omega \\ \omega \in Y, \omega' \notin Y}} p \cdot I(Y, 1, Bel(Y)) + (1 - p) \cdot I(Y, 0, Bel(Y)) \\ &\quad + \sum_{\substack{Z \subseteq \Omega \\ \omega' \in Z, \omega \notin Z}} (1 - p) \cdot I(Z, 1, Bel(Z)) + p \cdot I(Z, 0, Bel(Z)) . \end{aligned}$$

405 Now observe that every belief function $Bel^+ \in \mathbb{B}$ minimising $S_I^{aso}(P, \cdot)$ min-
 406 imises each of the four sums above individually, since every sum only depends
 407 on degrees of belief no other sum depends on.

408 By considering the first two sums for $U = \Omega$ and $W = \emptyset$ we find that
 409 $I(\Omega, 1, Bel^+(\Omega)) + I(\emptyset, 1, Bel^+(\emptyset))$ is uniquely minimised by $Bel^+(\Omega) = 1$
 410 and $Bel^+(\emptyset) = 0$.

411 Let us now consider the third sum. Note that any given $Y \subseteq \Omega$ such that
 412 $\omega \in Y$ and $\omega' \notin Y$ only appears in this sum once (and it does not appear
 413 in any other sum). Thus, $Bel^+(Y) = p = P(Y)$ is the unique minimum
 414 of $p \cdot I(Y, 1, \cdot) + (1 - p) \cdot I(Y, 0, \cdot)$. By varying $p = P(\omega)$ we obtain that
 415 $Bel^+(Y) = P(\omega)$ is the unique minimum of $p \cdot I(Y, 1, \cdot) + (1 - p) \cdot I(Y, 0, \cdot)$
 416 for all $p \in [0, 1]$ and all $Y \subseteq \Omega$ with $\omega \in Y$.

417 Finally, note that the above arguments do not depend on $\omega \in \Omega$. We thus
 418 find for all $Y \subseteq \Omega$ that $Bel^+(Y) = p$ is the unique minimum of $p \cdot I(Y, 1, \cdot) +$
 419 $(1 - p) \cdot I(Y, 0, \cdot)$ for all $p \in [0, 1]$.

420 Thus, IM_I is strictly proper. □

421 From a purely technical point of view, Theorem 5.6 can be most helpful.
 422 All one needs to do to check whether a SR S_I^{aso} is strictly \mathbb{B} -proper is to
 423 check whether the IM IM_I is strictly proper. The latter task can be accom-
 424 plished simply by checking whether simple sums are uniquely minimised by
 425 $Bel(X) = p$ and $Bel(\bar{X}) = 1 - p$. Checking strict \mathbb{B} -propriety requires one to
 426 solve a minimisation problem in $[0, 1]^{|\mathcal{P}\Omega|}$, which is in general a much harder
 427 problem.

428 Furthermore, Theorem 5.6 allows us to easily generate strictly \mathbb{B} -proper
 429 statistical SRs by association. That means that the class of inaccuracy mea-
 430 sures in our sense is a rich class consisting of a great variety of members.

431 We now give two applications of Theorem 5.6 in which we generate ex-
 432 tended strictly \mathbb{B} -proper SRs. The logarithmic IM ($I_{\log}(X, 1, x) := -\log(x)$,
 433 $I_{\log}(X, 0, x) := -\log(1 - x)$) and the spherical IM are well-known to be
 434 strictly proper ($I_{sph}(X, 1, x) := 1 + \frac{-x}{\sqrt{x^2 + (1-x)^2}}$, $I_{sph}(X, 0, x) := 1 + \frac{x-1}{\sqrt{x^2 + (1-x)^2}}$),
 435 see, e.g., [25, Section 8]).

436 **Corollary 5.7.** *The following logarithmic SR is strictly \mathbb{B} -proper.*

$$S_{\log}^{aso}(P, bel) := \sum_{X \subseteq \Omega} P(X) \cdot \left(-\log(Bel(X)) - \log(1 - Bel(\bar{X})) \right)$$

$$= \sum_{\omega \in \Omega} P(\omega) \cdot \left(- \sum_{\substack{X \subseteq \Omega \\ \omega \in \bar{X}}} \log(\text{Bel}(X)) - \log(1 - \text{Bel}(\bar{X})) \right) .$$

437 As usual in this context, we put $0 \cdot \infty := 0$ and $r \cdot \infty = \infty$ for $r \in (0, 1]$.
 438 By “log” we refer to a logarithm with an arbitrary base $b > 1$ and by
 439 “ln” to the natural logarithm, i.e., with base e .

440 **Corollary 5.8.** *The following spherical SR is strictly \mathbb{B} -proper.*

$$S_{sph}^{aso}(P, bel) := \sum_{X \subseteq \Omega} P(X) \cdot \left(2 + \frac{-\text{Bel}(X)}{\sqrt{\text{Bel}(X)^2 + (1 - \text{Bel}(X))^2}} + \frac{\text{Bel}(\bar{X}) - 1}{\sqrt{\text{bel}(\bar{X})^2 + (1 - \text{Bel}(\bar{X}))^2}} \right) .$$

441 For our running example, Brier Scores, we already considered the canon-
 442 ical association in Definition 5.2. We now note that Proposition 5.3 can
 443 alternatively be obtained as a simple corollary from Theorem 5.6
 444 Theorem 5.6 raises one, as of yet, open problem:

Open Problem 1: Does for all strictly \mathbb{B} -proper statistical SRs S_L exist an IM IM_I such that

$$S_L(P, Bel) = \sum_{\omega \in \Omega} P(\omega) \cdot IM_I(\omega, Bel) \quad ?$$

445 6. Justifying Probabilism with statistical Scoring Rules

446 In this section we build on Theorem 4.6 in order to obtain an epistemic
 447 justification of *Prob* for rational belief formation.

448 6.1. The Rationality of Tracking Objective Probabilities

449 We will assume the existence of objective probabilities and that the set
 450 of objective probability functions is \mathbb{P} . Whether such probabilities exist in
 451 the real world is a metaphysical debate, which we will not enter here. We
 452 content ourselves with noting that a number of writers have defended their
 453 existence in the real world. While the existence of objective probabilities in

454 the real world is a matter of debate, at least in (statistical) models featur-
 455 ing probability distributions of random variables objective probabilities may
 456 safely be assumed to exist.

457 Ideally, one might think, rational agents aim for beliefs which track the
 458 truth rather than tracking probabilities. Determining the truth, if such a
 459 thing as the true state of the world exists, has proven to be a rather com-
 460 plicated endeavour. Many have argued that if an agent knows the chances,
 461 then the only rational option is to set degrees of belief equal to the chances.
 462 We take it here that these arguments are right and that rational agents aim
 463 at tracking objective probabilities, at least in situations in which objective
 464 probabilities exist.

465 6.2. The formal Derivation

466 **Lemma 6.1.** *Let S_L be a strictly \mathbb{B} -proper SR. For all $Bel \in \mathbb{P}$ and all
 467 $Bel' \in \mathbb{B} \setminus \{Bel\}$ there exists a $P \in \mathbb{P}$ such that $S_L(P, Bel') > S_L(P, Bel)$.*

468 *Proof.* If $Bel \in \mathbb{P}$, then $S_L(Bel, \cdot)$ is uniquely minimized by $Bel = Bel$. So,
 469 for $Bel' \in \mathbb{B} \setminus \{Bel\}$ we have $S_L(Bel, Bel') > S_L(Bel, Bel)$. \square

470 **Theorem 6.2.** *Let S_I^{aso} be strictly proper and let IM_I be continuous.*

- 471 • For all $Bel \in \mathbb{B} \setminus \mathbb{P}$ there exists some $P \in \mathbb{P}$ such that for all $Q \in \mathbb{P}$
 472 $S_I^{aso}(Q, Bel) > S_I^{aso}(Q, P)$.
- 473 • For all $Bel \in \mathbb{P}$ and all $Bel' \in \mathbb{B} \setminus \{Bel\}$ there exists a $P \in \mathbb{P}$ such that
 474 $S_I^{aso}(P, Bel') > S_I^{aso}(P, Bel)$.

475 *Proof.* 1) Let $Bel \in \mathbb{B} \setminus \mathbb{P}$, then by Theorem 4.6 there exists a $P_{Bel} \in \mathbb{P}$ such
 476 that for all $\omega \in \Omega$ it holds that $IM_I(v_\omega, Bel) > IM_I(v_\omega, P_{Bel})$. For all $Q \in \mathbb{P}$
 477 there exists some $\omega \in \Omega$ such that $Q(\omega) > 0$. We thus find for all $Q \in \mathbb{P}$ that
 478 $S_I^{aso}(Q, Bel) > S_I^{aso}(Q, P_{Bel})$ holds.

479 2) By Theorem 5.6 S_I^{aso} is strictly \mathbb{B} -proper, now apply Lemma 6.1. \square

480 6.3. A brief Discussion

481 Besides the assumptions that rational agents aim only at accurate beliefs
 482 and that inaccuracy may be measured by a statistical SR S_L , the above justi-
 483 fication of *Prob* rests on the following: A) The statistical SR S_L is associated
 484 with an IM. B) S_I^{aso} is strictly \mathbb{B} -proper. C) Continuity of I . D) Dominance
 485 as DTN.

486 In order to make this justification compelling A – D need to be plausible.
 487 If rational agents only aim at accurate beliefs, then the statistical SR should
 488 be strictly \mathbb{B} -proper, as we argued in Section 3.2. If the answer to Open
 489 Problem 1 is “yes”, then B implies A. If the answer is “no”, then we either
 490 need to give an argument which singles out the class of statistical SRs which
 491 are associated with some IM IM_I or give a proof of Theorem 6.2 that also
 492 applies to statistical SRs which are not associated with an IM. Those who
 493 consider the class of strictly proper IMs to be the class inaccuracy measures in
 494 the epistemic approach seem to be forced to accept that the class of statistical
 495 SRs which measure inaccuracy by closeness-to-chances is precisely the class
 496 obtained by association.

497 Continuity is a fairly harmless technical condition. Again, as for A, it
 498 might be possible to prove Theorem 6.2 without assuming continuity.

499 As far as we are aware, no-one has seriously objected to dominance as
 500 DTN in this context, when applied to possible worlds. In the setting of
 501 this paper, agents aim at tracking objective probabilities and not at tracking
 502 worlds. It is thus fitting that dominance applies to objective probabilities in
 503 Theorem 6.2.

504 In Section 4.3 we argued that strict propriety for IMs without presup-
 505 posing that $Bel \in \mathbb{P}$ is unsatisfactory. For statistical SRs however, strict
 506 \mathbb{B} -propriety is desirable as a mean to encourage tracking of objective proba-
 507 bilities and thus reduce inaccuracy (Section 3.2). Under the assumption that
 508 strict propriety is technically necessary for convincing justifications of *Prob*,
 509 the upshot of Section 3.2 is that statistical SRs are *in principle* better suited
 510 than IMs for such justifications. Theorem 6.2 demonstrates that it is also
 511 *possible* to give a justification of *Prob* in the statistical framework.

512 The statistical approach has, at least in principle, one further advantage
 513 over the epistemic approach. Suppose the $\omega \in \Omega$ are the elementary events
 514 of some trial with chance distribution P^* . Given a belief function Bel and a
 515 SR S_L we can, at least in principle, approximate $S_L(P^*, Bel)$ by conducting
 516 i.i.d. trial runs. Thus, we do not need to have access to P^* to approximate
 517 $S_L(P^*, Bel)$. In the epistemic approach one assumes that there is an actual
 518 world ω^* among the $\omega \in \Omega$ but one does not know which possible world is
 519 the actual world. It is thus not possible, not even in principle, to compute
 520 $IM_I(v_{\omega^*}, Bel)$.

521 Another advantage distinct to the statistical approach is that it canoni-
 522 cally lends itself to take the agent’s evidence into account, as we shall see in
 523 the second part of this paper. The question of whether the classical epistemic

524 framework is able to adequately capture the agent’s evidence for justifications
525 of *Prob* is a matter of philosophical debate; see [13, 45]; which we will not
526 enter here.

527 6.4. Meeting some Objections

528 One may object that the here presented justification *presupposes* prob-
529 abilism by assuming the existence objective probability distributions which
530 satisfy Kolmogorov’s axioms. We openly acknowledge that we assumed the
531 existence of objective probabilities and that this assumption is key. Note how-
532 ever that the assumption of objective probabilities is an assumption about
533 the “outside world” which is *external* to the agent. We did not presuppose
534 anything about the agent’s degrees of belief (other than that they are real
535 numbers in $[0, 1] \subset \mathbb{R}$). Our presupposition thus concerns the agent’s envi-
536 ronment but not the agent’s doxastic state.

537 We want to make two further points. Firstly, justifications of *Prob* in the
538 framework of Section 4.2 which assume strict propriety presuppose *internal*
539 probabilism, the condition *strict propriety* involves an expectation! Secondly,
540 objective probabilities may well not exist in the real world. However, in
541 (toy) models their existence is guaranteed by the model specifications. The
542 sceptical reader may thus read our proposal as only applying to such toy
543 models. In general, we agree with Jaynes

544 In this connection we have to remember that probability theory
545 never solves problems of actual practice, because all such prob-
546 lems are infinitely complicated. We solve only idealizations of
547 the real problem, and the solution is useful to the extent that the
548 idealization is a good one. [23, p. 568]

549 One may also object that there are further epistemic goods which rational
550 agents ought to care for. It is certainly true that there might be other
551 epistemic goods, or even non-epistemic goods, rational agents ought to care
552 for. In the absence of a convincing account detailing what exactly these
553 goods are, we feel that it is appropriate to ignore these goods and solely
554 focus on inaccuracy minimisation.

555 The proponent of the classical epistemic framework in Section 4.2 may be
556 drawn to one of the following moves. Firstly, convincing justifications could
557 be given that do not require the IM IM_I to be strictly proper. This move
558 appears very unlikely, but possible, to succeed.

559 Secondly, one might head down the Joycean path and consider general
560 measures of inaccuracy $f(\omega, Bel)$. This path is, of course, open. The techni-
561 cal challenges one encounters appear to be so substantial, that assumptions
562 need to be made which make the justifications less than fully satisfactory.

563 Thirdly, an argument may be advanced claiming that the class of appro-
564 priate IMs is a proper subclass of the strictly proper IMs. The appeal of
565 such an approach then hinges on the characterisation of this subclass of IMs.
566 Such an argument was put forward in [30, 31]. The class of IMs considered in
567 [30, 31] is so narrow that it does not contain the logarithmic nor the spher-
568 ical IM. Their justification, improving on de Finetti’s result by moderately
569 enlarging the class of IMs, can thus only be a step towards a satisfactory
570 justification of *Prob*. Until such a reasonably large subclass of strictly proper
571 IMs has been discovered, we remain sceptical about this approach.

572 Part 2

573 7. Maximum Entropy Principles

574 The first part of this paper focussed on justifications of *Prob*. A great
575 number of writers invoke further norms to constrain the choice of a belief
576 function more tightly. Typically, such norms are Calibration Norms ([63,
577 Section 3.3]), a Principal Principle ([33, 43, 47]) or the Maximum Entropy
578 Principle (discussed in more detail below) to constrain the choice of a belief
579 function depending on the agent’s evidence. Justifications of such approaches
580 are normally given in a two-stage argument. First, *Prob* is justified, then the
581 further norm(s) are justified. This leaves proponents of such approaches
582 with the complicated task of explaining why and how the justification of
583 *Prob* supersedes the justification(s) of the further invoked norm(s).

584 In this section we give a single justification for *Prob* and Maximum Gen-
585 eralised Entropy Principles at the same time. Since we give a single justifi-
586 cation no two-stage justificatory argument is required of the proponent of a
587 combination of *Prob* and a Maximum Generalised Entropy Principle.

588 Exactly as in the first part, we do not presuppose *Prob*, strict \mathbb{P} -propriety
589 is hence of little use. The key notion will again be strict \mathbb{B} -propriety.

590 As in the first part of this paper we focus on formal aspects of the justi-
591 fications and only touch on the question as to when DTNs apply. The DTN
592 we will here use is Worst-Case Expected Loss (WCEL) avoidance. In the for-
593 mal literature, WCEL has rich history and goes back to the seminal work of
594 Morgenstern and von Neumann. The most obvious toy cases in which WCEL
595 avoidance is an appropriate DTN are two-player single-round games with an
596 adversary playing after Player1 has made her move. Recently, normative ar-
597 guments for risk sensitivity were advanced in [4]. A maximally risk-sensitive
598 agents adheres to WCEL avoidance.

599 The justifications we give here apply to interpretations of $P \in \mathbb{E}$ as epis-
600 temic subjective probabilities or as objective probabilities.

601 7.1. The general Arguments

602 Consider an agent with current evidence which narrows the chance func-
603 tion down to a non-empty and convex set $\emptyset \subset \mathbb{E} \subseteq \mathbb{P}$. \mathbb{E} is called the set of
604 *calibrated* functions. The most prominent objective Bayesian approach then
605 requires an agent to equivocate sufficiently between the basic propositions

606 that the agent can express while adopting a belief function in \mathbb{E} , cf. [63].⁵
 607 That is, the agent is required to assign the basic propositions the same prob-
 608 abilities as far as this is consistent with the agent’s evidence. This norm is
 609 then spelled out in terms of the Maximum Entropy Principle:

Maximum Entropy Principle (MaxEnt) A rational agent ought to adopt a probability function $Bel \in \mathbb{E}$ which maximises Shannon Entropy, H_{\log}

$$H_{\log}(Bel) := \sum_{\omega \in \Omega} -Bel(\omega) \log(Bel(\omega)) . \quad (20)$$

610 The probability function $P_{=} \in \mathbb{P}$ defined by $P_{=}(\omega) := \frac{1}{|\Omega|}$ for all $\omega \in \Omega$ is
 611 called the *equivocator*. $P_{=}$ is the function in \mathbb{P} with greatest entropy. MaxEnt
 612 can be understood as requiring an agent to adopt a belief function in \mathbb{E} which
 613 is as similar to $P_{=}$ as possible.

614 MaxEnt has given rise to a substantial literature on rational belief for-
 615 mation; as examples we mention [1, 8, 23, 29, 40, 41].

616 Key to MaxEnt is the loss function $L(\omega, Bel) = -\log(Bel(\omega))$ and the
 617 logarithmic scoring rule S_{\log}

$$S_{\log}(P, Bel) := \sum_{\omega \in \Omega} -P(\omega) \log(Bel(\omega)) .$$

618 We can express Shannon Entropy in terms of this SR, $H_{\log}(P) = S_{\log}(P, P)$.

619 MaxEnt is well-known to be justified on the following grounds of WCEL
 620 avoidance [20, 59]

621 **Theorem 7.1** (Justification of MaxEnt). *If $\emptyset \neq \mathbb{E} \subseteq \mathbb{P}$ is convex and closed,*
 622 *then*

$$\arg \inf_{Bel \in \mathbb{P}} \sup_{P \in \mathbb{E}} S_{\log}(P, Bel) = \arg \sup_{P \in \mathbb{E}} H_{\log}(P) \quad (21)$$

⁵For our purposes, it is not relevant to explain what “sufficiently equivocates” amounts to. We shall only be concerned with maximal equivocation.

623 and there is only one unique such function maximising Shannon Entropy,
 624 P^\dagger .

625 So, an agent which aims to minimise $\sup_{P \in \mathbb{E}} S_L(P, Bel)$ by adopting a
 626 probabilistic belief function $Bel \in \mathbb{P}$, i.e., avoiding worst-case expected loga-
 627 rithmic loss, has to adopt P^\dagger as her belief function.

628 We now generalise this well-known justification of MaxEnt to strictly \mathbb{X} -
 629 proper SRs which satisfy the following minimax equation

$$\inf_{Bel \in \mathbb{X}} \sup_{P \in \mathbb{E}} S_L(P, Bel) = \sup_{P \in \mathbb{E}} \inf_{Bel \in \mathbb{X}} S_L(P, Bel) . \quad (22)$$

630 See [58] for an introduction to such minimax equations which arose from Von
 631 Neumann's seminal game theoretical work.

632 Following [20], we call $H_L(P) := S_L(P, P)$ *generalised entropy*. If the set
 633 $\arg \sup_{P \in \mathbb{E}} H_L(P)$ contains a unique function, then this function is denoted
 634 by P^\ddagger and called *generalised entropy maximiser*. The following generalises
 635 [20, Theorem 6.4] to non-probabilistic belief functions.

636 **Theorem 7.2** (Justification of Generalised Entropy Maximisation). *If $\emptyset \neq$
 637 $\mathbb{E} \subseteq \mathbb{P}$ is convex and closed, S_L strictly \mathbb{X} -proper, (22) holds and if $H_L(P)$ is
 638 strictly concave on \mathbb{P} , then*

$$\arg \inf_{Bel \in \mathbb{X}} \sup_{P \in \mathbb{E}} S_L(P, Bel) = \arg \sup_{P \in \mathbb{E}} H_L(P) =: \{P^\ddagger\} . \quad (23)$$

639 *Proof.* Let us first use (22) and the fact that S_L is strictly \mathbb{X} -proper to obtain

$$\inf_{Bel \in \mathbb{X}} \sup_{P \in \mathbb{E}} S_L(P, Bel) = \sup_{P \in \mathbb{E}} \inf_{Bel \in \mathbb{X}} S_L(P, Bel) \quad (24)$$

$$= \sup_{P \in \mathbb{E}} S_L(P, P) . \quad (25)$$

640 Since \mathbb{E} is convex, closed and non-empty the function $S_L(P, P)$ has a unique
 641 supremum in \mathbb{E} . That is, the set $\arg \sup_{P \in \mathbb{E}} \sum_{\omega \in \Omega} S_L(P, P)$ consists of a
 642 unique probability function which is in \mathbb{E} , P^\ddagger .

643 Using \mathbb{X} -strict propriety to obtain the strict inequality in (27) we find for
 644 all $Bel \in \mathbb{X} \setminus \{P^\ddagger\}$

$$\sup_{P \in \mathbb{E}} S_L(P, Bel) \geq S_L(P^\ddagger, Bel) \quad (26)$$

$$> S_L(P^\ddagger, P^\ddagger) . \quad (27)$$

645 Recall that $\inf_{Bel \in \mathbb{X}} \sup_{P \in \mathbb{E}} S_L(P, Bel)$ equals $S_L(P^\ddagger, P^\ddagger)$. Thus, no $Bel \in$
 646 $\mathbb{X} \setminus \{P^\ddagger\}$ minimises $\sup_{P \in \mathbb{E}} S_L(P, Bel)$. Hence, P^\ddagger is the unique minimiser of
 647 $\sup_{P \in \mathbb{E}} S_L(P, Bel)$. \square

648 This means that an agent which aims to minimise $\sup_{P \in \mathbb{E}} S_L(P, Bel)$ by
 649 adopting a belief function $Bel \in \mathbb{B}(!)$, i.e., avoiding worst-case expected loss,
 650 has to adopt P^\ddagger as her belief function.

651 So, Theorem 7.2 simultaneously justifies *Prob* and the following principle:

652 **Maximum Generalised Entropy Principle** A rational agent
 653 ought to adopt the unique probability function in \mathbb{E} which max-
 654 imises the generalised entropy $H_L(P)$.

655 The question arises how P^\ddagger changes when the agent receives new infor-
 656 mation and the set of calibrated functions changes. It is not rational for
 657 a WCEL avoiding agent to change her belief, if $\mathbb{E}' \subset \mathbb{E}$ and $P^\ddagger \in \mathbb{E}'$ (see
 658 below). This property of unchanged beliefs has been termed *obstinacy*, see
 659 for example [40, p. 80].

660 **Corollary 7.3.** *Let \mathbb{E} and S_L be as in Theorem 7.2. If $\emptyset \subset \mathbb{E}' \subset \mathbb{E}$ contains*
 661 *P^\ddagger , then*

$$\arg \inf_{Bel \in \mathbb{X}} \sup_{P \in \mathbb{E}'} S_L(P, Bel) = \{P^\ddagger\} . \quad (28)$$

662 Note that we do not require that \mathbb{E}' is convex nor that \mathbb{E}' is closed.

663 *Proof.* First note that

$$\inf_{Bel \in \mathbb{X}} \sup_{P \in \mathbb{E}'} S_L(P, Bel) \leq \inf_{Bel \in \mathbb{X}} \sup_{P \in \mathbb{E}} S_L(P, Bel) \quad (29)$$

$$= S_L(P^\ddagger, P^\ddagger) . \quad (30)$$

664 For all $Bel \in \mathbb{X} \setminus \{P^\ddagger\}$ we find using strict \mathbb{X} -propriety of S_L that

$$\sup_{P \in \mathbb{E}'} S_L(P, Bel) \geq S_L(P^\ddagger, Bel) \quad (31)$$

$$> S_L(P^\ddagger, P^\ddagger) . \quad (32)$$

665 For the belief function P^\ddagger we find

$$S_L(P^\ddagger, P^\ddagger) \leq \sup_{P \in \mathbb{E}'} S_L(P, P^\ddagger) \quad (33)$$

$$\leq \sup_{P \in \mathbb{E}} S_L(P, P^\ddagger) \quad (34)$$

$$= S_L(P^\ddagger, P^\ddagger) . \quad (35)$$

666 So, $\sup_{P \in \mathbb{E}'} S_L(P, P^\ddagger) = S_L(P^\ddagger, P^\ddagger)$. Hence, P^\ddagger uniquely minimises WCEL.
667 □

668 7.2. Generalised Entropies

669 Theorem 7.2 gives general conditions under which generalised entropy
670 maximisation is justified with respect to the choice of a *particular* statistical
671 SR. Unsurprisingly, the choice of different SRs, i.e., utility functions, leads
672 to different generalised entropy maximisers. The importance of choosing an
673 appropriate SR has recently been emphasised in [36].

674 Consider the extended Brier score S_{Brier}^{ext} , the spherical SR S_{sph}^{aso} and
675 $S_{llog}^{ext} := -\frac{|P\Omega|}{2} + \sum_{Y \subseteq \Omega} Bel(Y) - \sum_{X \subseteq \Omega} P(X) \cdot \ln(Bel(X))$. We now show
676 that all three SRs satisfy the conditions in Theorem 7.2. We shall
677 not give the rather uninformative calculations but rather state the result of
678 these calculations.

679 All three SRs are strictly \mathbb{B} -proper, see Proposition 5.3, Corollary 5.8 and
680 Proposition 9.1.

681 Straightforward calculations show that Brier Entropy $H_{Brier}(P)$ and the
682 Spherical Entropy $H_{Sph}(P)$ are strictly concave on \mathbb{P} . The entropy of the
683 logarithmic SR is $H_{P\Omega}(P) := \sum_{X \subseteq \Omega} -P(X) \log(P(X))$ which we shall prove
684 in Section 9.1. This entropy is called *Proposition Entropy* in [29]. Clearly,
685 $H_{P\Omega}$ is strictly concave on \mathbb{P} .

686 Note that $H_{P\Omega}$ is different from Shannon Entropy, H_{log} . In $H_{P\Omega}$ the sum
687 is taken over all events $X \subseteq \Omega$ and not over all elementary events $\omega \in \Omega$. Not
688 only are Proposition Entropy and Shannon Entropy different functions; in

689 general, their respective maximum obtains for different probability functions
 690 in \mathbb{E} , cf. [29, Figure 1, p. 3536].

691 That all three entropies considered here are sufficiently regular, satisfying
 692 the minimax condition (24), follows for instance from König’s result [28, p.
 693 56], see [51] for a discussion of König’s result.

694 These three entropies have different maximisers on rather simple sets \mathbb{E} ,
 695 as can be gleaned from Figure 1 and Figure 2.

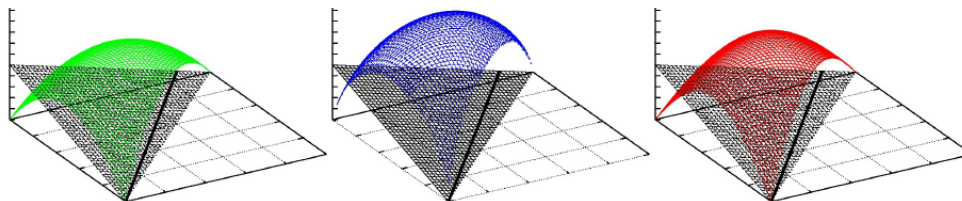


Figure 1: Brier Entropy H_{Brier} (green), Proposition Entropy $H_{P\Omega}$ (blue) and Spherical Entropy H_{Sph} (red) for $\Omega = \{\omega_1, \omega_2, \omega_3\}$. The black line segment connects $P_1 = (1, 0, 0)$ and $P_2 = (0, \frac{5}{6}, \frac{1}{6})$.

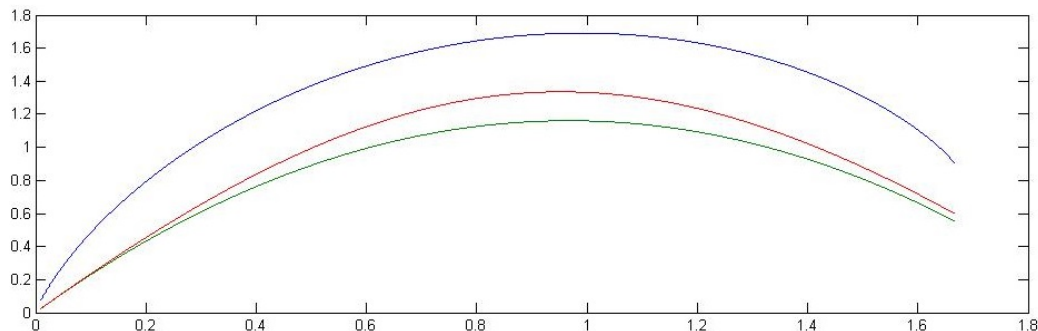


Figure 2: Brier Entropy H_{Brier} (green), Proposition Entropy $H_{P\Omega}$ (blue) and Spherical Entropy H_{Sph} (red) plotted along the line segment between $P_1 = (1, 0, 0)$ and $P_2 = (0, \frac{5}{6}, \frac{1}{6})$ parametrised as $P_1 + t \cdot (-0.6, 0.5, 0.1)$ for $t \in [0, \frac{10}{6}]$. The Brier Entropy maximiser is $P_{Brier}^\dagger = (0.4194, 0.4839, 0.0968)$ [$t = 0.968$], the Proposition Entropy maximiser is $P_{P\Omega}^\dagger = (0.4054, 0.4955, 0.0991)$ [$t = 0.991$] and the Spherical Entropy maximiser is $P_{Sph}^\dagger = (0.4277, 0.4770, 0.0954)$ [$t = 0.954$]. The absolute value of the Spherical Entropy has been adjusted to fit all curves neatly into the picture.

696 Theorem 7.2 deals with generalised entropies. The question arises whether
 697 we can find a statistical SR to simultaneously justify *Prob* and *MaxEnt*. Un-
 698 fortunately, we do not know the answer to this question

Open Problem 2 Does there exist a strictly \mathbb{B} -proper statistical SR S_L such that (24) holds and such that for all closed and convex $\emptyset \subset \mathbb{E} \subseteq \mathbb{P}$ it holds that

$$\arg \sup_{P \in \mathbb{E}} S_L(P, P) = \arg \sup_{P \in \mathbb{E}} H_{\log}(P) ? \quad (36)$$

699 *7.3. Generalised Entropies and the Principle of Indifference*

700 The Principle of Indifference (PoI) has long fascinated philosophers. We
 701 here show that maximising generalised entropies entails the PoI for many
 702 natural generalised entropies. Recent arguments in its favor can be found in
 703 [37, 46, 62].

704 **Definition 7.4.** A SR S_L is called equivocator neutral, if and only if for all
 705 $\omega, \omega' \in \Omega$ it holds that $L(\omega, P_{=}) = L(\omega', P_{=})$.

706 **Theorem 7.5** (Generalised Entropies and PoI). *If S_L is equivocator neutral,*
 707 *strictly \mathbb{X} -proper with $\mathbb{P} \subseteq \mathbb{X} \subseteq \mathbb{B}$, satisfies (24) and if $H_L(P)$ is strictly*
 708 *concave on \mathbb{P} , then*

$$\arg \inf_{Bel \in \mathbb{X}} \sup_{P \in \mathbb{P}} S_L(P, Bel) = \arg \sup_{P \in \mathbb{P}} S_L(P, P) = \{P_{=}\} . \quad (37)$$

709 So, under complete ignorance, $\mathbb{E} = \mathbb{P}$, the unique rational choice under
 710 WCEL avoidance is $Bel = P_{=}$; this provides a justification of the PoI. For a
 711 recent justification of the PoI using IMs we refer the reader to [46].

712 *Proof.* From Theorem 7.2 and the fact that \mathbb{P} is convex and closed we obtain

$$\arg \inf_{Bel \in \mathbb{X}} \sup_{P \in \mathbb{P}} S_L(P, Bel) = \arg \sup_{P \in \mathbb{P}} S_L(P, P) . \quad (38)$$

713 Note that since S_L is equivocator neutral, there exists some constant
 714 $c \in \mathbb{R}$ such that for all $\omega \in \Omega$ it holds that $L(\omega, P_{=}) = c$.

715 Assume for contradiction that there exists some $Q \in \arg \sup_{P \in \mathbb{P}} H_L(P)$
 716 which is different from $P_{=}$. Since $H_L(P)$ is a strictly concave function on \mathbb{P}
 717 the maximum of $H_L(\cdot)$ has to be unique and hence $H_L(Q) > H_L(P)$. We
 718 then obtain using (38)

$$H_L(P_=) < H_L(Q) \quad (39)$$

$$= \inf_{Bel \in \mathbb{X}} \sup_{P \in \mathbb{P}} S_L(P, Bel) \quad (40)$$

$$\leq \sup_{P \in \mathbb{P}} S_L(P, P_=) \quad (41)$$

$$= \sup_{P \in \mathbb{P}} \sum_{\omega \in \Omega} P(\omega) L(\omega, P_=) \quad (42)$$

$$= \sup_{P \in \mathbb{P}} \sum_{\omega \in \Omega} P(\omega) c \quad (43)$$

$$= c \quad (44)$$

$$= \sum_{\omega \in \Omega} \frac{1}{|\Omega|} L(\omega, P_=) \quad (45)$$

$$= \sum_{\omega \in \Omega} P_=(\omega) L(\omega, P_=) \quad (46)$$

$$= H_L(P_=) . \quad (47)$$

719 Contradiction. Thus, $\{P_=\} = \arg \sup_{P \in \mathbb{P}} H_L(P)$. □

720 Equivocator neutrality is a very weak symmetry condition on L . Strict
 721 \mathbb{B} -propriety and satisfying (24) are standing assumptions in this section. Fi-
 722 nally, $\arg \sup_{P \in \mathbb{P}} H_L(P)$ containing a unique element would follow from H_L
 723 being strictly concave. If S_L is strictly \mathbb{P} -proper, then H_L is concave, see [17,
 724 p. 361]. Thus, in a large number of cases maximising generalised entropy
 725 entails the PoI.

726 Not only is the equivocator the unique function minimising WCEL under
 727 complete ignorance, it is also the unique such function as long as $P_= \in \mathbb{E}$:

728 **Corollary 7.6.** *For a SR S_L as in Theorem 7.5 and for all sets $\mathbb{E} \subset \mathbb{P}$ such*
 729 *that $P_= \in \mathbb{E}$ it holds that*

$$\arg \inf_{Bel \in \mathbb{X}} \sup_{P \in \mathbb{E}} S_L(P, Bel) = \arg \sup_{P \in \mathbb{P}} S_L(P, P) = \{P_=\} . \quad (48)$$

730 *Proof.* First, let us reason as in Theorem 7.5 to obtain the equality below

$$\inf_{Bel \in \mathbb{X}} \sup_{P \in \mathbb{E}} S_L(P, Bel) \leq \inf_{Bel \in \mathbb{X}} \sup_{P \in \mathbb{P}} S_L(P, Bel) \quad (49)$$

$$= S_L(P_-, P_-) . \quad (50)$$

731 Using strict propriety we find for all belief functions $Bel \in \mathbb{X} \setminus \{P_-\}$ that

$$\sup_{P \in \mathbb{E}} S_L(P, Bel) \geq S_L(P_-, Bel) \quad (51)$$

$$> S_L(P_-, P_-) . \quad (52)$$

732 So all belief functions different from the equivocator P_- have a strictly sub-
 733 optimal WCEL. P_- has the best possible WCEL as we saw in Theorem 7.5.
 734 It follows that

$$\arg \inf_{Bel \in \mathbb{X}} \sup_{P \in \mathbb{E}} S_L(P, Bel) = \{P_-\} . \quad (53)$$

735 □

736 For instance, $S_L = S_{Brier}^{ext}$, $S_L = S_{sph}^{aso}$ and $S_L = S_{llog}^{ext}$ satisfy the assump-
 737 tions of Theorem 7.5. We hence obtain

738 **Corollary 7.7.** *If $S_L = S_{Brier}^{ext}$, $S_L = S_{sph}^{aso}$ or $S_L = S_{llog}^{ext}$ and $P_- \in \mathbb{E}$, then*

$$\arg \inf_{Bel \in \mathbb{B}} \sup_{P \in \mathbb{E}} S_L(P, Bel) = \arg \sup_{P \in \mathbb{P}} S_L(P, P) = \{P_-\} . \quad (54)$$

739 8. Local Scoring Rules

740 We now turn our attention to strictly \mathbb{B} -proper statistical SRs themselves.
 741 S_{log} stands out as the only strictly \mathbb{P} -proper local SR and as the heart of Max-
 742 Ent. It has hence received considerable attention in the literature. Locality
 743 means that if a elementary event $\omega \in \Omega$ obtains, then the loss incurred only
 744 depends on $Bel(\omega)$ and not on the entire belief function.

745 Subsequently, we will take an interest in notions of locality applied to
 746 SRs defined on $\mathbb{P} \times \mathbb{B}$. Surprisingly, the most natural way of extending the
 747 notion of locality to $\mathbb{P} \times \mathbb{B}$ is incompatible with strict \mathbb{B} -propriety.

748 8.1. Locality and strict \mathbb{P} -propriety

749 **Definition 8.1.** A statistical SR $S_L : \mathbb{P} \times \mathbb{P} \rightarrow [0, +\infty]$ is called local, if and
 750 only if $L(\omega, Bel)$ only depends on the belief in ω and not on other beliefs.
 751 Abusing the notation in the usual way we write $L(Bel(\omega))$.

752 The class of such SRs which are strictly \mathbb{P} -proper is rather simple:

753 **Theorem 8.2** (Savage [54]). *Up to an affine-linear transformation, the only*
 754 *local and strictly \mathbb{P} -proper statistical SR is*

$$S_{\log}(P, Bel) = \sum_{\omega \in \Omega} -P(\omega) \log(Bel(\omega)) . \quad (55)$$

755 Local SRs or logarithmic loss functions have been argued for in a variety
 756 of settings. For example, in [66, pp. 16] and [2, p. 72-73] for belief elicitation.
 757 See [7, p. 2039-2040] for a discussion on locality and [7, p. 2046] for an
 758 axiomatic characterisation of logarithmic SRs in terms of scale-invariance.

759 Levinstein points out advantages of S_{\log} as a measure of inaccuracy over
 760 S_{Brier} applied to probabilistic belief functions, see [32]. We also want to
 761 mention that S_{\log} is the only strictly \mathbb{P} -proper SR which is consistent with
 762 the use of likelihoods or log likelihoods to evaluate assessors, cf. [64, p. 1075].
 763 In [63, p. 64-65], Williamson shows that S_{\log} can be characterised in terms of
 764 four natural axioms, one of which is locality. S_{\log} has found applications in a
 765 variety of areas, for example in information theory [9, 52], Neyman-Pearson
 766 Theory in statistics [15] and the health sciences [26].

767 Recently, the IM_{\log} has left a positive impression in formal epistemology
 768 as a tool to measure a degree of confirmation, see [60].

769 Let us now consider a local loss function $L : [0, 1] \rightarrow [0, +\infty]$ and the
 770 corresponding local SR $S_L : \mathbb{P} \times \mathbb{B} \rightarrow [0, +\infty]$ (defined on belief functions
 771 $Bel \in \mathbb{B}$!)

$$S_L(P, Bel) = \sum_{\omega \in \Omega} P(\omega) \cdot L(Bel(\omega)) . \quad (56)$$

772 Note that only beliefs in elementary events appear in the above expression.
 773 Thus, beliefs in non-elementary events will not affect the score $S_L(P, Bel)$.
 774 Thus, a DTN applying local statistical SR $S_L(P, Bel)$ can only yield con-
 775 straints on the agent's beliefs in elementary events; beliefs in non-elementary

776 events are completely unconstrained. So, local SRs are ill-suited for justifi-
 777 cations of norms of rational belief formation without presupposing *Prob*.

778 Thus, we now investigate how to extend the notion of locality, which
 779 proved to be technically fruitful when *Prob* was presupposed, without pre-
 780 supposing *Prob*.

781 8.2. Locality, strict \mathbb{B} -propriety and extended Scoring Rules

782 One obvious way to generalise locality is:

783 **Definition 8.3.** *An extended SR is called ex-local, if and only if there exists*
 784 *a loss function $L_{loc} : \mathcal{P}\Omega \times [0, 1] \rightarrow [0, \infty]$ such that*

$$S_{L_{loc}}^{ext}(P, Bel) = \sum_{X \subseteq \Omega} P(X) \cdot L_{loc}(X, Bel(X)) \quad (57)$$

$$= \sum_{\omega \in \Omega} P(\omega) \cdot \left(\sum_{\substack{X \subseteq \Omega \\ \omega \in X}} L_{loc}(X, Bel(X)) \right) . \quad (58)$$

785 Ex-locality here means that $L(X, Bel)$ is of the form $L_{loc}(X, Bel(X))$, i.e.
 786 the loss attributable to event X in isolation of all other events, if X obtains,
 787 only depends on X and on $Bel(X)$. Here, we do *not* allow $L(X, Bel(X))$ to
 788 depend on further beliefs such as $Bel(\bar{X})$.

789 This notion of an ex-local extended SR differs from *local* statistical SRs in
 790 Savage's sense in two respects. Firstly, the sum is now over all events $X \subseteq \Omega$
 791 and not only over the elementary events $\omega \in \Omega$. Secondly, the loss function
 792 L_{loc} may now depend on the event X whereas Savage's loss function only
 793 depended on the belief in an elementary event ω and not in the elementary
 794 event itself.

795 If S_L^{ext} is ex-local, then the loss attributable to $Bel(X)$ only enters once
 796 into (57). More precisely, the only summand depending on $Bel(X)$ is $P(X) \cdot$
 797 $L_{loc}(X, Bel(X))$. Since P is a probability function, $P(\emptyset) = 0$ holds. Hence,
 798 by our convention that $0 \cdot \infty = 0$ we obtain $P(\emptyset) \cdot L_{loc}(\emptyset, Bel(\emptyset)) = 0 \cdot$
 799 $L_{loc}(\emptyset, Bel(\emptyset)) = 0$ for all $P \in \mathbb{P}$. So, $S_{L_{loc}}^{ext}(P, Bel)$ does not depend on
 800 $Bel(\emptyset)$.

801 Hence, a belief function Bel_a which agrees with P on all events $\emptyset \subset X \subseteq \Omega$
 802 and $Bel(\emptyset) = a$ with $a \in (0, 1]$ it holds that $S_{L_{loc}}^{ext}(P, P) = S_{L_{loc}}^{ext}(P, Bel_a)$.
 803 Thus, no ex-local SR is strictly \mathbb{B} -proper.

804 One might initially think that the incompatibility of ex-locality and strict
805 \mathbb{B} -propriety is only due to the fact that for all $P \in \mathbb{P}$ $P(\emptyset) = 0$ holds.
806 However, we shall now see that this is not the case.

807 Let $\mathbb{B}^- := \{Bel : \mathcal{P}\Omega \setminus \{\emptyset\} \rightarrow [0, 1]\}$ and define strict \mathbb{B}^- -propriety of a SR
808 S_L in the obvious way, i.e., for all $P \in \mathbb{P}$ it holds that $\arg \inf_{Bel \in \mathbb{B}^-} S(P, Bel) =$
809 $\{P|_{\mathcal{P}\Omega \setminus \{\emptyset\}}\}$. For ease of notation we drop the restriction operator “ $|$ ” from
810 now on.

811 **Theorem 8.4.** *There does not exist an ex-local extended strictly \mathbb{B}^- -proper*
812 *SR $S_{L_{loc}}^{ext}$.*

813 *Proof.* It is sufficient to show that for all $P \in \mathbb{P}$

$$\arg \inf_{Bel \in \mathbb{B}^-} S_{L_{loc}}^{ext}(P, Bel) = \arg \inf_{Bel \in \mathbb{B}^-} \sum_{X \subseteq \Omega} P(X) \cdot L_{loc}(X, Bel(X)) \quad (59)$$

814 does not depend on P , since strict \mathbb{B}^- -propriety would require that the above
815 minimum obtains uniquely for $Bel = P$.

816 For a fixed loss function L_{loc} and a fixed event $\emptyset \subset X \subseteq \Omega$ it holds that
817 $\arg \inf_{Bel(X) \in [0,1]} L_{loc}(X, Bel(X))$ only depends on $Bel(X) \in [0, 1]$ and not on
818 P nor on $Bel(Y)$ for $Y \neq X$. Furthermore, $Bel(X)$ may be freely chosen in
819 $[0, 1]$, since Bel does not have to satisfy any further constraints, such as the
820 axioms of probability. Hence, for all $\emptyset \subset X \subseteq \Omega$ the infimum (or infima) of
821 $P(X)L_{loc}(X, Bel(X))$ obtains independently of P .

822 Thus, $S_{L_{loc}}^{ext}(P, Bel)$ is minimised, if and only if every summand in (59)
823 is minimised. For each summand this minimum obtains independently of
824 P . □

825 **Proposition 8.5.** $S_{\log}^{ext}(P, Bel) := \sum_{X \subseteq \Omega} -P(X) \cdot \log(Bel(X))$ is not strictly
826 \mathbb{B}^- -proper.

827 *Proof.* Define a belief function $Bel_1 \in \mathbb{B}$ by $Bel_1(X) := 1$ for all $X \subseteq \Omega$. For
828 all $P \in \mathbb{P}$ and all $X \subseteq \Omega$ it holds that $P(X) \log(Bel_1(X)) = 0$. So, for all
829 $P \in \mathbb{P}$

$$Bel_1 \in \arg \inf_{Bel \in \mathbb{B}} S_{\log}^{ext}(P, Bel) . \quad (60)$$

830 □

831 Recall from Theorem 8.2 that the logarithmic SR S_{\log} is the only local
 832 \mathbb{P} -strictly proper statistical SR. Evidently, strict propriety crucially depends
 833 on the set of scored belief functions.

834 The SR considered in Corollary 5.7: $S_{\log}^{aso}(P, Bel) := \sum_{X \subseteq \Omega} P(X) \cdot (-\log(Bel(X)) -$
 835 $\log(1 - Bel(\bar{X})))$ is not ex-local. The loss term depends on $Bel(X)$ and
 836 $Bel(\bar{X})$. Thus, Proposition 5.7 does not contradict Theorem 8.4.

837 Note that $S_{Loc}^{ext}(P, Bel)$ does not depend on $Bel(X)$ for all those event
 838 $X \subset \Omega$ with $P(X) = 0$. If any genuine measure of inaccuracy has to take
 839 into account how $P(X)$ and $Bel(X)$ relate for *all* $X \subseteq \Omega$, then ex-local
 840 SRs cannot serve as measures of inaccuracy. In this case, the impossibility
 841 theorem only rules out the existence of SRs which are unsuitable for our
 842 purposes.

843 9. Two Notions of Locality

844 The question we now pose is: how much of the locality condition do we
 845 need to give up in order obtain strictly \mathbb{B} -proper extended SRs which are
 846 local, *in some sense*?

847 9.1. Penalties

848 As it turns out, there exists an extended SR employing logarithms which
 849 is strictly \mathbb{B} -proper.

850 **Proposition 9.1.** *The following extended SR is strictly \mathbb{B} -proper*

$$851 \quad S_{ll\log}^{ext}(P, Bel) := \sum_{X \subseteq \Omega} P(X) \cdot \left(-1 + \frac{\sum_{Y \subseteq \Omega} Bel(Y)}{\sum_{Y \subseteq \Omega} P(Y)} - \ln(Bel(X)) \right) \quad (61)$$

$$852 \quad = -\frac{|\mathcal{P}\Omega|}{2} + \sum_{Y \subseteq \Omega} Bel(Y) - \sum_{X \subseteq \Omega} P(X) \cdot \ln(Bel(X)) . \quad (62)$$

851 This SR is not purely logarithmic since it contains the *penalty term*,
 852 $\sum_{Y \subseteq \Omega} Bel(Y)$. This term penalises belief functions for indiscriminately as-
 853 signing high degrees of belief to all events. In particular it prevents $Bel_1 \in \mathbb{B}$
 854 from being the score minimiser. The penalty term is constant for all $X \subseteq \Omega$,
 855 it is thus global.

Proof. Define an IM $IM_{ll\log}$ by

$$\begin{aligned} I(X, 0, Bel(X)) &:= Bel(X) \\ I(X, 1, Bel(X)) &:= Bel(X) - 1 - \ln(Bel(X)) . \end{aligned}$$

856 We now show that $IM_{ll\log}$ is strictly proper. Clearly, $IM_{ll\log}$ is never strictly
857 less than zero.

858 Let $p \in [0, 1]$ and $\emptyset \subset X \subset \Omega$ be fixed and let

$$\begin{aligned} f(Bel(X)) &:= p \cdot I(X, 1, Bel(X)) + (1 - p) \cdot I(X, 0, Bel(X)) \\ &= p \cdot Bel(X) - p - p \cdot \ln(Bel(X)) + (1 - p) \cdot Bel(X) \\ &= -p - p \cdot \ln(Bel(X)) + Bel(X) . \end{aligned}$$

859 By equating the derivative of $f(Bel(X))$ with zero we find for $p > 0$

$$\frac{df(Bel(X))}{dBel(X)} = -\frac{p}{Bel(X)} + 1 = 0 . \quad (63)$$

860 Trivially, this equation is uniquely solved by $Bel(X) = p > 0$. Considering
861 the second derivative of $f(Bel(X))$ shows that $Bel(X) = p > 0$ is the unique
862 minimum.

863 For $p = 0$ we recall the usual convention that $0 \ln(Bel(X)) = 0$, even if
864 $Bel(X) = 0$. Hence, $f(Bel(X)) = (1 - p) \cdot I(X, 0, Bel(X)) = Bel(X)$, which
865 is uniquely minimised by $Bel(X) = p = 0$.

For $X = \emptyset$ and $X = \Omega$ we have

$$I_{ll\log}(\Omega, 1, Bel(X)) + I_{ll\log}(\emptyset, 0, Bel(X)) = Bel(\Omega) - 1 - \ln(Bel(\Omega)) - Bel(\emptyset),$$

866 which is uniquely minimised by $Bel(\Omega) = 1$ and $Bel(\emptyset) = 0$.

867 We next show that $S_{ll\log}^{ext}$ is strictly \mathbb{B} -proper. We do so by showing that
868 it is associated with $IM_{ll\log}$ and hence strictly \mathbb{B} -proper by Theorem 5.6.

$$\sum_{\omega \in \Omega} P(\omega) \cdot IM_{ll\log}(\omega, Bel)$$

$$\begin{aligned}
&= \sum_{\omega \in \Omega} P(\omega) \cdot \left(\sum_{\substack{X \subseteq \Omega \\ \omega \in X}} I(X, 1, Bel(X)) + \sum_{\substack{Y \subseteq \Omega \\ \omega \notin Y}} I(Y, 0, Bel(Y)) \right) \\
&= \sum_{\omega \in \Omega} P(\omega) \cdot \left(\sum_{\substack{X \subseteq \Omega \\ \omega \in X}} Bel(X) - 1 - \ln(Bel(X)) + \sum_{\substack{Y \subseteq \Omega \\ \omega \notin Y}} Bel(Y) \right) \\
&= \sum_{\omega \in \Omega} P(\omega) \cdot \left(\sum_{Z \subseteq \Omega} Bel(Z) + \sum_{\substack{X \subseteq \Omega \\ \omega \in X}} -1 - \ln(Bel(X)) \right) \\
&= \sum_{Z \subseteq \Omega} Bel(Z) + \sum_{\omega \in \Omega} P(\omega) \cdot \left(\sum_{\substack{X \subseteq \Omega \\ \omega \in X}} -1 - \ln(Bel(X)) \right) \\
&= \sum_{Z \subseteq \Omega} Bel(Z) + \sum_{X \subseteq \Omega} P(X) \cdot \left(-1 - \ln(Bel(X)) \right) \\
&= \sum_{X \subseteq \Omega} P(X) \cdot \sum_{Z \subseteq \Omega} \frac{Bel(Z)}{\sum_{Y \subseteq \Omega} P(Y)} + \sum_{X \subseteq \Omega} P(X) \cdot \left(-1 - \ln(Bel(X)) \right) \\
&= \sum_{X \subseteq \Omega} P(X) \cdot \left(\frac{\sum_{Z \subseteq \Omega} Bel(Z)}{\sum_{Y \subseteq \Omega} P(Y)} - 1 - \ln(Bel(X)) \right) \\
&= S_{llog}^{ext}(P, Bel) .
\end{aligned}$$

869

□

870 S_{llog}^{ext} contains a local term, $\ln(Bel(X))$, and a global term, $\sum_{Y \subseteq \Omega} Bel(Y)$.
871 The constant term, $-\frac{|P\Omega|}{2}$, has been added for the following cosmetic reason.
872 For $Bel \in \mathbb{P}$ we have

$$S_{llog}^{ext}(P, Bel) = - \sum_{X \subseteq \Omega} P(X) \cdot \ln(Bel(X)) \quad (64)$$

$$= S_{log}^{ext}(P, Bel) . \quad (65)$$

873 So, for $Bel \in \mathbb{P}$ we recapture the SR considered in Proposition 8.5 (for the
874 natural logarithm) and we note that

$$S_{llog}^{ext}(P, P) = - \sum_{X \subseteq \Omega} P(X) \cdot \ln(P(X)) .$$

875 At first glance, S_{llog}^{ext} appears to be an extended strictly \mathbb{B} -proper SR which
876 is not associated to an IM. If this were the case, then we would have solved
877 Open Problem 1 (Section 5.2) in the negative. However, we saw in the above
878 proof that S_{llog}^{ext} is indeed associated with the strictly proper IM_{llog} . We have
879 thus not solved Open Problem 1.

880 Finally, let us remark that proving strict \mathbb{B} -propriety of S_{llog}^{ext} directly is a
881 rather complicated endeavour. The above proof is a nice illustration of the
882 technical helpfulness of Theorem 5.6 to which we alluded to in Section 5.2.

883 9.2. Normalising Beliefs

884 In Proposition 9.1 we saw how one can use a penalty term to construct a
885 strictly \mathbb{B} -proper logarithmic SR. In [29] the authors showed that the penalty
886 term can be dropped, if the belief functions are normalised, that is the belief
887 functions considered are in some set $\mathbb{B}_{norm} \supset \mathbb{P}$.

888 We shall now quickly summarise the relevant points in [29]: Denote by
889 π a set of non-empty mutually exclusive, jointly exhaustive proper subsets
890 of Ω , i.e., a partition. Denote by Π the union of $\{\Omega, \emptyset\}$, $\{\Omega\}$ and the set of
891 these partitions. Then define

$$\mathbb{B}_{norm} := \{B : \mathcal{P}\Omega \rightarrow [0, 1] \mid \sum_{F \in \pi} B(F) = 1 \text{ for some } \pi \in \Pi$$

$$\text{and } \sum_{F \in \pi} B(F) \leq 1 \text{ for all } \pi \in \Pi\} .$$

892 For a given a weighting function $g : \Pi \rightarrow \mathbb{R}_{\geq 0}$ such that for all $\emptyset \subseteq X \subseteq \Omega$ it
893 holds that $\sum_{\substack{\pi \in \Pi \\ X \in \pi}} g(\pi) > 0$, a SR is defined on $\mathbb{P} \times \mathbb{B}_{norm}$ by:

$$S_{normlog,g}^{ext}(P, B) := - \sum_{\pi \in \Pi} g(\pi) \sum_{X \in \pi} P(X) \cdot \log(B(X)) \quad (66)$$

$$= \sum_{X \subseteq \Omega} P(X) \cdot \left(\sum_{\substack{\pi \in \Pi \\ X \in \pi}} g(\pi) \right) \cdot \log(B(X)) . \quad (67)$$

894 **Proposition 9.2.** [29, Corollary 3, p. 3542] $S_{normlog,g}^{ext}(P, B)$ is strictly \mathbb{B}_{norm} -
895 proper for all such g .

896 Note that since $\mathbb{P} \subset \mathbb{B}_{norm}$, strict \mathbb{B}_{norm} -propriety is well defined in the
897 sense of Definition 3.2.

898 The above proposition does not contradict Theorem 8.4, since we here
899 consider normalised belief functions in \mathbb{B}_{norm} while Theorem 8.4 concerns
900 belief functions in \mathbb{B} .

901 The SRs S_{llog}^{ext} and $S_{normlog,g}^{ext}$ rely on the same idea: The main culprit in
902 the impossibility Theorem 8.4 is that in (59) there is no interaction between
903 the degrees of belief in different events. Normalising beliefs re-introduces
904 such an interaction. The main structural difference between the two SRs is
905 how normalisation is achieved. The former SR, S_{llog}^{ext} , introduces a penalty
906 (i.e. normalisation) term into the SR, for the latter SR, $S_{normlog,g}^{ext}$, one pre-
907 supposes normalised belief functions.

908 10. Conclusion

909 In the first part of this paper we saw how to use statistical SRs to justify
910 *Prob*. In this second part we demonstrated the usefulness of statistical SRs for
911 justifications of further norms of rational belief formation. In particular, we
912 saw how an agent’s evidence can be naturally taken into account by applying
913 WCEL avoidance as DTN.

914 Logarithmic SRs occupy a prominent place in the literature as protago-
915 nists in Savage’s theorem and objective Bayesianism. We hence set out to
916 investigate how to construct statistical logarithmic SRs which are strictly
917 \mathbb{B} -proper. We found three such logarithmic SRs (Proposition 5.7, Proposi-
918 tion 9.1 and Proposition 9.2).

919 Ideas from the epistemic and the statistical approach have been influential
920 in the development of this paper. Looking into the future, pulling strands
921 from both approaches together appears to have the potential to be benefi-
922 cial for both approaches. Generally speaking, extending Richard Pettigrew’s
923 Epistemic Utility Theory Programme [42, 48] to statistical SRs appears to
924 be a research avenue holding great promise. We thus hope for many more
925 exciting entries to be added to Table 1.

926 Unfortunately, we did not answer all the questions we raised. Hopefully,
927 future work will solve the problems left open in this paper.

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Decision Theoretic Norm	Inaccuracy Measures	Scoring Rules
Dominance w.r.t. $\omega \in \Omega$	[11], [49], [24], [25], [43], [44]	[55] [56]
Dominance w.r.t. $P \in \mathbb{P}$		Corollary 5.4, Theorem 6.2
Expected Loss w.r.t. Bel^*		Belief Elicitation
Worst-Case Loss w.r.t. $\omega \in \Omega$	[46]	
WCEL w.r.t. $P \in \mathbb{E}$		Theorems 7.1, 7.2, 7.5 [20], [29]

Table 1: Combinations of IMs and SRs with DTNs

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